## Formal Techniques – 2015-09-03

**Exercise 1.** Let  $f : A \to B$  be an arbitrary function between the DCPOs A, B. Assume that, for any directed  $D \subseteq A$ , we have that  $\bigsqcup^B f[D]$  (exists and) is equal to  $f(\lvert \mid^A D)$ . Then, prove that f is monotonic.

**Exercise 2.** Consider the following protocol excerpt written in the applied-pi notation.

Apply the control flow analysis to the protocol above, generating a tree automaton to over-approximate the message flow, as done by function gen(...). Provide a list of states for such automaton and the transitions among them. For each state, briefly hint to its relationship with the protocol above.

Exercise 3. Consider the following tree automaton

 $\begin{array}{l} @a: \cos(@c, @b), \cos(@g, @f), \operatorname{enc}(@d, @e), \operatorname{dec}(@a, @a).\\ @b: \operatorname{fst}(@a).\\ @c: \operatorname{snd}(@a).\\ @d: \operatorname{m.}\\ @e: \cos(@f, @g).\\ @f: k1.\\ @g: k2. \end{array}$ 

and the rewriting rules

 $\operatorname{dec}(\operatorname{enc}(M,K),K) \Rightarrow M \qquad \operatorname{fst}(\operatorname{cons}(X,Y)) \Rightarrow X \qquad \operatorname{snd}(\operatorname{cons}(X,Y)) \Rightarrow Y$ 

Apply the completion algorithm to the above automaton, building an over-approximation for the languages associated to its states. Assuming @a models the set of messages being exchanged over a public channel, state what can be concluded about the secrecy of message m.

**Exercise 4.** Formally prove the following formula exploiting the Curry-Howard isomorphism.

$$\forall p, q : \mathsf{Prop.} ((p \to q) \lor q) \to (p \to q)$$

**Exercise 5.** For each  $i \in \{1, 2\}$ , let  $C_i$  and  $A_i$  be CLs with a Galois connection  $\alpha_i : C_i \xrightarrow{\leftarrow} A_i : \gamma_i$ . Construct a Galois connection

$$\alpha : [\mathcal{C}_1 \to \mathcal{C}_2] \stackrel{\leftarrow}{\to} [\mathcal{A}_1 \to \mathcal{A}_2] : \gamma$$

where  $[- \rightarrow -]$  denotes the CL of Scott-continuous functions. Prove that yours is indeed a Galois connection.

**Exercise 6.** Let A be a CL, and  $f : A \to A$  be a monotonic function. Prove that  $fix(f) = fix(f \circ f)$ , where fix denotes the minimum fixed point.